Storing EPICS process variables in HDF5 files for ITER

Rodrigo Castro¹, Yury Makushok², Lana Abadie³, Bertrand Bauvir³, Ralph Lange³, Andre Neto⁴

> ¹Laboratorio Nacional de Fusión, CIEMAT. Madrid. Spain ²MINSAIT - INDRA, Madrid, Spain ³ITER Organization, St Paul lez Durance, France ⁴ Fusion for Energy, Barcelona, Spain

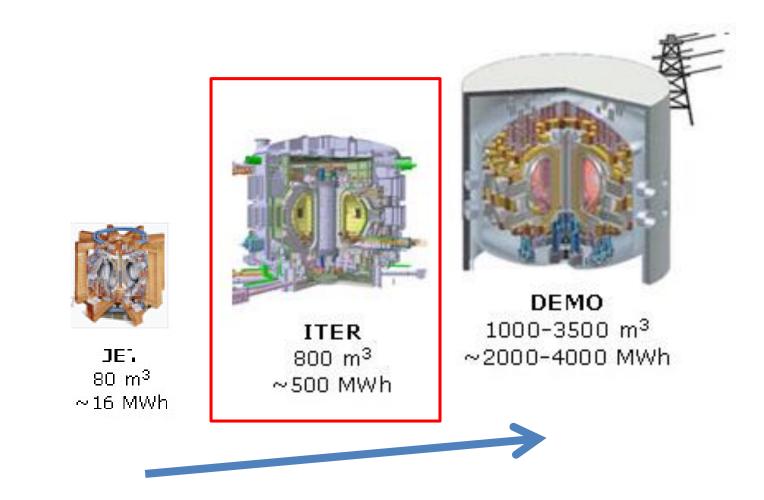
> > 2022 European HDF5 User Group 01/06/2022







Current context



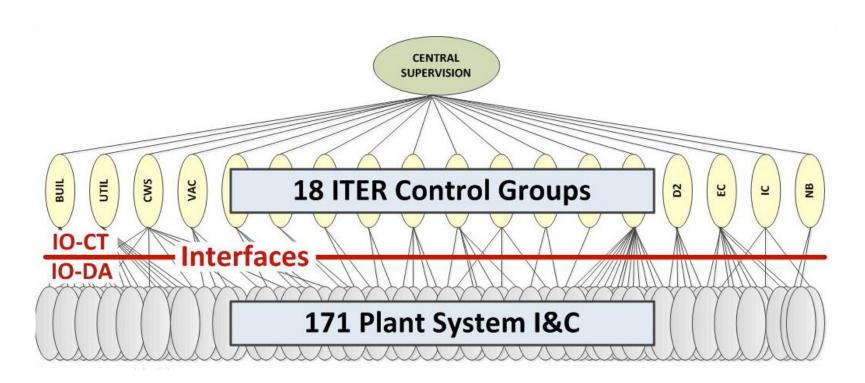








ITER control system organization



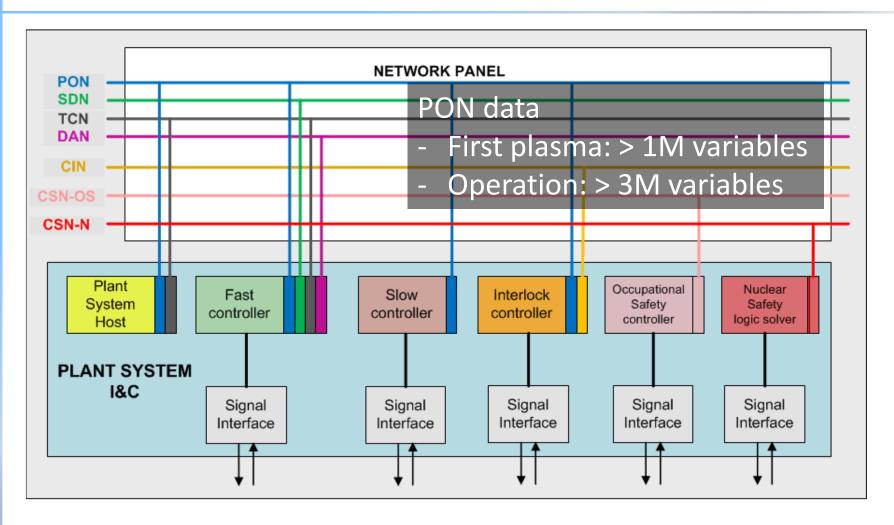
ITER CODAC: Common language for all PS I&C
 Distributed control system based on EPICS







Plant Operation Network (PON) data









Design and implementation consideration

Time evolution data in steady-state operation

- Requests based on time interval
- Different files along the time can be involved
- Different data nature along the time can be involved
 - Different dimensionality
 - Different sampling rate
 - Different units
 - Different datatype

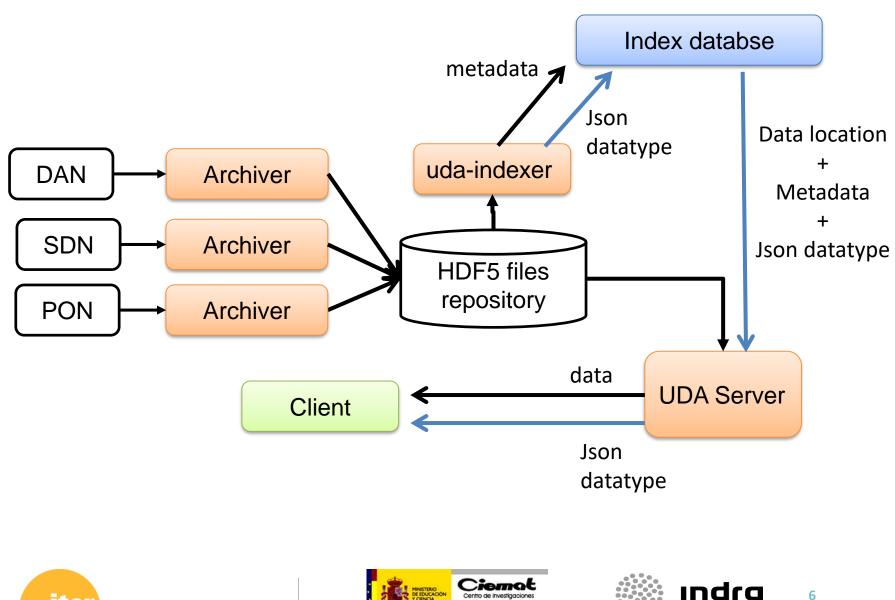
Data must be accessible on fly

- Data flush timeouts
- Real time indexing mechanism
 - New files / growing files
 - New files, open for writing files and closed for writing files must be notified





Data archiving/retrieving cycle





6

EPICS Process Variables

Before EPICS v7: Channel Access Protocol

Values have a simple composite type

- Timestamp
- Status
- Severity
- Value
- Support a set of primitive data types: integer, double, string, enum
- Some metadata can be retrieved
 - Enum labels
 - Visualization metadata







EPICS Process Variables

*

□ After EPICS v7: PVAccess protocol

- Very complex nested datatypes can be defined
- Can include multiplicity at any level
- Structure leafs have additional metadata

Example

Root		
	ain	
	SDNHeader	
	Time	
-	Sensor[32]	
	 Integrated 	
	State	int32
	Quality	int32
	Value	float32
	Error	float32
	Proportional	
	IntegratedFiltered	
	ProportionalFiltered	
	Combined	
	CombinedIntegrated	
	Temperature	
	ErrorFlags	uint32
	FPGAInfo	
	▼ FPGAMon	
	✓ Vols	2022/02/2010/04/04 2
	PLInternal	float32
	PLAuxil[4]	float32
	PLBlockRAM	float32
	PSLowPowerDomain	float32
	PSAuxil[4]	float32
	Temperature	float32
	TemperatureAlarmCount	uint32
	 VoltAIC 	
	PLInternal	uint32
	PLAuxil[4]	uint32
	PLBlockRAM	uint32
	PSLowPowerDomain	uint32
	PSAuxil[4]	uint32
	ClockErrorsCount	uint32
	ClockFrequencies[6]	float32
	InternalErrorsCount	uint32
	ConfigHash	uint32





Channel Access archiving: PON archiver

✓ <u>SWMR model</u>

- Aggregates the maximum number of PVs in one file
- Implements flush timeout (to warranty maximum read latency)
- HDF5 backend files rotate in some size / time conditions
- Any PV can change on fly:
 - 🗸 Туре
 - Enum labels
 - 🗸 Display metadata





PON archiver (Channel Access archiver)

HDF5 model

- One group per PV
 - Payload dataset: timestamp, status, severity, value
 - Dynamic metadata attributes
 - Enum labels
 - Display metadata
- Until 80K PVs / file





PON archiver (Channel Access archiver)

□ Challenges

- Memory problem
 - HDF5 caches must be correctly managed
- Expensive file rotation
 - Creation of all objects have big CPU usage and takes significant time
 - File rotation in asyn mode (separate thread)
- If a change in PV property: datatype, units, enum labels
 - New individual file is created until next rotation
- Performance problem for flushing updated PVs
 - One flush for all file at the end of the loop (avoiding flush per dataset)





PVaccess archiving: PVA archiver

✓ <u>SWMR model</u>

- Manages big nested datatypes
- Autodiscovery PV datatype
- ✓ PV datatype can change





Big nested datatype: some alternatives

- Use opac datatype + datatype definition
 - Breaks our current model based on HDF5 types
 - Mandatory to read all data structure just for one field
- □ Flat nested structure: 1 composite field
 - Cases of 16K fields (10 DAQ boards)
 - HDF5 limitations: maximum about 1300 fields
- One dataset per field
 - Good read performance
 - Poor write performance: 1 write -> 16K writes





Big nested datatype: first implementation

Flat structure break algorithm

- □ Iterates flatten structure trying to:
 - Find the longest common path (trying to group as much as possible)
 - Until aggregation limit (number of fields or size limit) is achieved
 - Check if this path already exists in the file
 - If exists -> necessary to force a new break with 1 level longer paths (less aggregation)
 - Add a group for the found path name
 - With a composite fields payload dataset that aggregates all data under the found path name





Big nested datatype: first implementation

Flat structure break algorithm

Pros:

- Current HDF5 archiving model (reading, indexing) is valid
- Level of aggregation (size of datasets) is configurable
- Good aggregation results / universal algorithm : Breaks 8K fields (5 DAQ boards) -> 220 datasets

Cons:

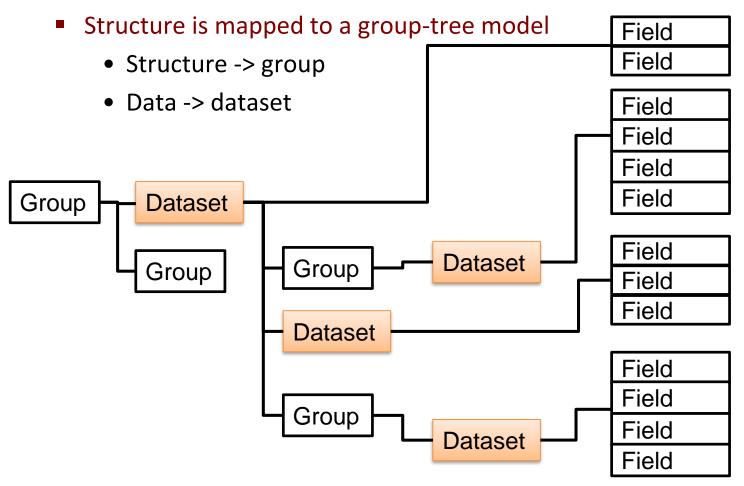
 HDF5 structure is not visually a 1-to-1 map of the original nested structure





Big nested datatype: second implementation

Group-tree data model







Big nested datatype: second implementation

Group-tree data model

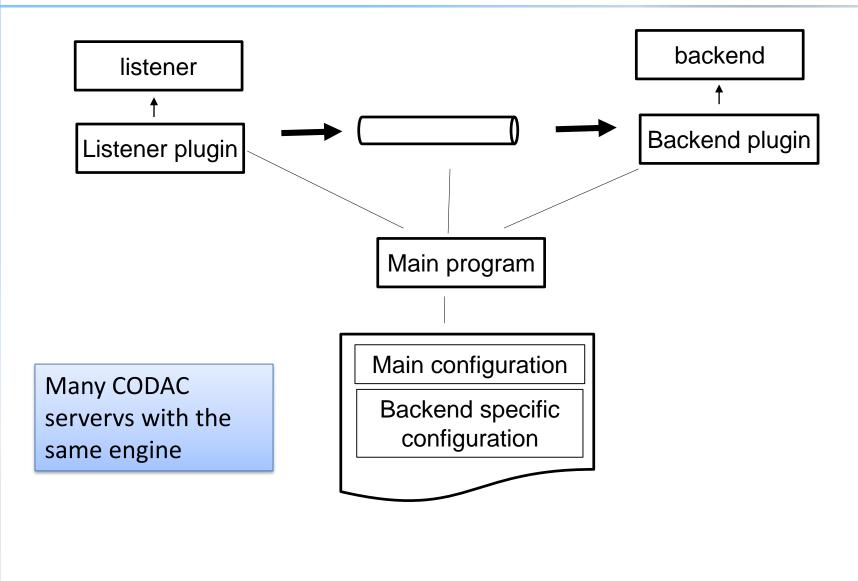
Pros:

- HDF5 structure maps 1-to-1 the original nested structure
- Easy to extract a subtree of data
- Read performance
- Cons:
 - Not already a complete solution in case of leaf composite datatype with more than 1300 fields (is really a limitation?)





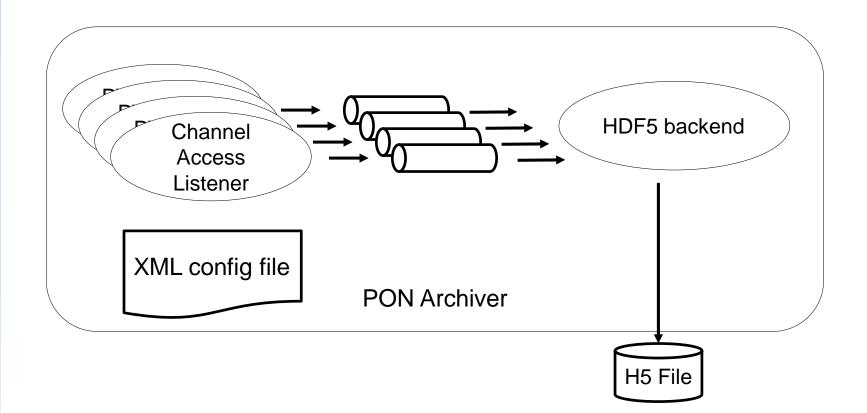
Common Architecture







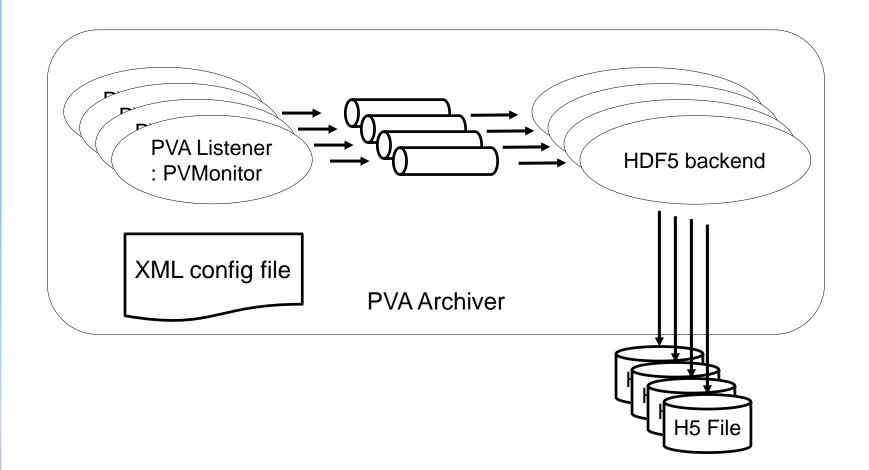
PON archiver implementation







PVA archiver implementation









Thank you for your attention

Questions?





