Predicting and optimizing the performance of HDF5 applications

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Software stack

Previous modeling work



Applications directly access data. Previous works build models to predict the I/O time.

How to predict the end-to-end performance of the array I/O requests?

• This work



Predict

Motivation

• Read an entire array



Optimize

Read path in HDF5



The cache read and memory copy operations can spend up to 90% and 87% of the execution time in our experiments.

Predict

An end-to-end model should cover all the steps in the I/O path.

End-to-end time = cache read time + data server access time + memory copy time

Solution overview



- The system parameters
 - the chunk shape
 - storage format (HDF5 or Zarr)
 - stripe count and size in the file system
 - cache size
- The solution contains three components
 - The cache state component maintains the status of the library cache and the file system cache

Predict

- The variable calculation component computes the variables in the model based on the selected region
- The model component predicts the cache read time, data server access time and memory copy time

Model accuracy



Real and predicted time

Experiment variables

Predict

Variable	Value range
Array shape	$36M \times 64$, $2.4K \times 2K \times 3K$
Selected region	entire array, 1K rows, 1 column, 1K boxes
Chunk length in Dim 1	70K - 4.5M
Chunk length in Dim 2	1 - 64
Number of processes	1 - 64
Stripe count	1 - 64
File system	Lustre, GPFS

- Takeaways:
 - The RMSE of the model is 0.29
 - The model correctly predicts the fastest library between HDF5 and Zarr 94% of the time.

• Supernova detection



21×21 pixels/image

Optimize

Vortices prediction



Optimize

- One I/O per array
 - I/O takes 200× to 700× longer than computation





Predict

Too many small I/O requests!

Optimize

HDF5 on heterogeneous data stores



Optimize

- Store arrays on heterogeneous data stores
 - Without modifying applications
- Accelerate small I/O requests
 - Placement \rightarrow improve the performance of one request
 - Consolidation \rightarrow reduce the number of requests
- Automatically decide the array storage layout
 - Which data store should an array be placed in?
 - How do we store small arrays in chunks?

Optimize

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Optimize

System architecture



Optimize

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Optimize

I/O acceleration techniques

• Placement



Consolidation



Predict



Chunk 1 Chunk 2

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Optimize

Optimization workflow Predict



Optimization impact



Directly read many small arrays from HadoopFS

Predict

Consolidate small arrays with same access pattern in a chunk. Read in fewer I/O requests.

Place hot arrays and co-accessed arrays in Redis. Read less data from HadoopFS.

Conclusions

- Reading from cache and memory copy to transform layouts can spend 90% of the execution time
 - An end-to-end model to cover the entire I/O path
- Applications spend 99% of the time to read small arrays
 - VOL transparently forward requests to two data stores
 - Placement and consolidation reduce the number of I/O requests to the slow data store
 - An analytical cost model helps to decide the storage layout